## PAC-Bayesian Learning of Optimization Algorithms



Peter Ochs<br>Mathematical Optimization for Data Science Saarland University<br>- 07.09.2023 -<br>

joint work: Michael Sucker

## Example 1:

$$
\min _{x} f(x), \quad f(x):=\frac{1}{2}\|A x-b\|^{2} .
$$

## Example 1:

$$
\min _{x} f(x), \quad f(x):=\frac{1}{2}\|A x-b\|^{2} .
$$

How do we solve the problem?

## Example 1:

$$
\min _{x} f(x), \quad f(x):=\frac{1}{2}\|A x-b\|^{2} .
$$

How do we solve the problem?
$\diamond$ Inspect the properties of the problem.
Example: Smooth/Quadratic problem with $L=\|A\|^{2}$-Lipschitz gradient.

## Inverse Problems are often Modelled as an Optimization Problem

## Example 1:

$$
\min _{x} f(x), \quad f(x):=\frac{1}{2}\|A x-b\|^{2}
$$

How do we solve the problem?
$\diamond$ Inspect the properties of the problem.
Example: Smooth/Quadratic problem with $L=\|A\|^{2}$-Lipschitz gradient.
Embed the problem into a class of problems for which algorithms are available.
Example: Use Gradient Descent with step size $\alpha=1 / L$

$$
x^{(k+1)}=x^{(k)}-\alpha \nabla f\left(x^{(k)}\right)
$$

Worst case convergence guarantee:

$$
f\left(x^{(k)}\right)-\min f \leq O(1 / k)
$$

## Hidden Structures

## If we knew ...

## If we knew ...

that $A$ is actually of the form

$$
A=\left(\begin{array}{cccc}
10 & 0 & \cdots & 0 \\
0 & 1 & \ddots & \vdots \\
\vdots & \ddots & \ddots & 0 \\
0 & \cdots & 0 & 1
\end{array}\right) ?
$$

## Hidden Structures

If we knew ...
$\diamond$ that $A$ is actually of the form

$$
A=\left(\begin{array}{cccc}
10 & 0 & \cdots & 0 \\
0 & 1 & \ddots & \vdots \\
\vdots & \ddots & \ddots & 0 \\
0 & \cdots & 0 & 1
\end{array}\right) \quad ?
$$

$\diamond$ We would write down a different algorithm (that directly returns the solution).
$\diamond$ Game-changer, if many such problems for different $b$ need to be solved.
$\diamond$ Sometimes the "best" class of problems is not obvious!

## Hidden Structures

If we knew ...
$\diamond$ that $A$ is actually of the form

$$
A=\left(\begin{array}{cccc}
10 & 0 & \cdots & 0 \\
0 & 1 & \ddots & \vdots \\
\vdots & \ddots & \ddots & 0 \\
0 & \cdots & 0 & 1
\end{array}\right) \quad ?
$$

$\diamond$ We would write down a different algorithm (that directly returns the solution).
$\diamond$ Game-changer, if many such problems for different $b$ need to be solved.
$\diamond$ Sometimes the "best" class of problems is not obvious!

Can we construct an algorithm that adapts to hidden problem structures?

## Example with Noisy Operator $A$

Example 2: Solve many problems of the form

$$
\min _{x} f_{A}(x), \quad f_{A}(x):=\frac{1}{2}\|A x-b\|^{2} \quad \text { where } \quad A=\bar{A}+\text { noise } .
$$

## Example with Noisy Operator $A$

Example 2: Solve many problems of the form

$$
\min _{x} f_{A}(x), \quad f_{A}(x):=\frac{1}{2}\|A x-b\|^{2} \quad \text { where } \quad A=\bar{A}+\text { noise } .
$$

## Using Gradient Descent:

$\diamond$ For each problem $f_{A}$, we need to compute $L=\|A\|^{2}$,
$\diamond$ and run Gradient Descent with $\alpha=1 / L$ to solve the problem.
$\diamond$ Computation of $\|A\|$ can be expensive.

## Example with Noisy Operator $A$

Example 2: Solve many problems of the form

$$
\min _{x} f_{A}(x), \quad f_{A}(x):=\frac{1}{2}\|A x-b\|^{2} \quad \text { where } \quad A=\bar{A}+\text { noise } .
$$

## Using Gradient Descent:

$\diamond$ For each problem $f_{A}$, we need to compute $L=\|A\|^{2}$,
$\diamond$ and run Gradient Descent with $\alpha=1 / L$ to solve the problem.
$\diamond$ Computation of $\|A\|$ can be expensive.

Or ...
$\diamond$ if the noise is bounded, we can use a worst case estimate for $L$.
$\diamond$ Results in small step sizes.
$\diamond$ Upper bound may be too pessimistic for most problems in practice.

## Example with Noisy Operator $A$

Example 2: Solve many problems of the form

$$
\min _{x} f_{A}(x), \quad f_{A}(x):=\frac{1}{2}\|A x-b\|^{2} \quad \text { where } \quad A=\bar{A}+\text { noise } .
$$

## Using Gradient Descent:

$\diamond$ For each problem $f_{A}$, we need to compute $L=\|A\|^{2}$,
$\diamond$ and run Gradient Descent with $\alpha=1 / L$ to solve the problem.
$\diamond$ Computation of $\|A\|$ can be expensive.

Or ...
$\diamond$ if the noise is bounded, we can use a worst case estimate for $L$.
$\diamond$ Results in small step sizes.
$\diamond$ Upper bound may be too pessimistic for most problems in practice.

Can we construct an algorithm with good performance for more likely problems?

## Data Driven Approach

Yes, using data driven approaches / learning !

## Data Driven Approach

Learning alleviates the bounds of analytical tractability by providing more:
$\diamond$ Information: Leverage more structure.
Automation: Less "hand-crafting".
Possibilities: More building blocks.

## Data Driven Approach

## Yes, using data driven approaches / learning !

Learning alleviates the bounds of analytical tractability by providing more:
$\diamond$ Information: Leverage more structure.
Automation: Less "hand-crafting".
Possibilities: More building blocks.

## Our goals:

$\diamond$ Breaking the barrier of worst-case estimates.
$\diamond$ Adapt algorithms to hidden problem structures.
$\diamond$ Define tight classes of problems.
$\diamond$ We insist on having some theoretical guarantees.

Consider the random parametric optimization problem:

$$
\min _{x \in \mathbb{R}^{n}} \ell(x, \mathfrak{S})
$$

$\diamond \ell: \mathbb{R}^{n} \times \Theta \rightarrow \mathbb{R}_{\geq 0}$ is a given measurable loss-function.
$\diamond \mathfrak{S}:(\Omega, \mathcal{F}, \mathbb{P}) \rightarrow \Theta$ is a random variable.

Consider the random parametric optimization problem:

$$
\min _{x \in \mathbb{R}^{n}} \ell(x, \mathfrak{S})
$$

$\diamond \ell: \mathbb{R}^{n} \times \Theta \rightarrow \mathbb{R}_{\geq 0}$ is a given measurable loss-function.
$\diamond \mathfrak{S}:(\Omega, \mathcal{F}, \mathbb{P}) \rightarrow \Theta$ is a random variable.

## Example:

Regularized Inverse Problem:

$$
\ell(x, \lambda)=\frac{1}{2}\|A x-b\|^{2}+\lambda R(x), \quad \text { i.e. } \theta:=\lambda, \Theta=[0,1]
$$

Consider the random parametric optimization problem:

$$
\min _{x \in \mathbb{R}^{n}} \ell(x, \mathfrak{S})
$$

$\diamond \ell: \mathbb{R}^{n} \times \Theta \rightarrow \mathbb{R}_{\geq 0}$ is a given measurable loss-function.
$\diamond \mathfrak{S}:(\Omega, \mathcal{F}, \mathbb{P}) \rightarrow \Theta$ is a random variable.
$\diamond$ Use a parametric optimization algorithm, i.e. a measurable function:

$$
\mathcal{A}: \mathcal{H} \times \mathbb{R}^{n} \times \Theta \rightarrow \mathbb{R}^{n}, \quad\left(\alpha, x^{(0)}, \theta\right) \mapsto \mathcal{A}\left(\alpha, x^{(0)}, \theta\right)
$$

## Example:

Regularized Inverse Problem:

$$
\ell(x, \lambda)=\frac{1}{2}\|A x-b\|^{2}+\lambda R(x), \quad \text { i.e. } \theta:=\lambda, \Theta=[0,1] .
$$

Consider the random parametric optimization problem:

$$
\min _{x \in \mathbb{R}^{n}} \ell(x, \mathfrak{S})
$$

$\diamond \ell: \mathbb{R}^{n} \times \Theta \rightarrow \mathbb{R}_{\geq 0}$ is a given measurable loss-function.
$\diamond \mathfrak{S}:(\Omega, \mathcal{F}, \mathbb{P}) \rightarrow \Theta$ is a random variable.
$\diamond$ Use a parametric optimization algorithm, i.e. a measurable function:

$$
\mathcal{A}: \mathcal{H} \times \mathbb{R}^{n} \times \Theta \rightarrow \mathbb{R}^{n}, \quad\left(\alpha, x^{(0)}, \theta\right) \mapsto \mathcal{A}\left(\alpha, x^{(0)}, \theta\right)
$$

## Example:

Regularized Inverse Problem:

$$
\ell(x, \lambda)=\frac{1}{2}\|A x-b\|^{2}+\lambda R(x), \quad \text { i.e. } \theta:=\lambda, \Theta=[0,1] .
$$

Concatenation of a fixed number of Preconditioned Gradient Descent steps:

$$
x^{(k+1)}=x^{(k)}-P \nabla \ell\left(x^{(k)}, \theta\right), \quad \text { i.e. } \alpha:=P, \mathcal{H}:=\mathbb{R}^{n \times n} .
$$

Consider the random parametric optimization problem:

$$
\min _{x \in \mathbb{R}^{n}} \ell(x, \mathfrak{S})
$$

$\diamond \ell: \mathbb{R}^{n} \times \Theta \rightarrow \mathbb{R}_{\geq 0}$ is a given measurable loss-function.
$\diamond \mathfrak{S}:(\Omega, \mathcal{F}, \mathbb{P}) \rightarrow \Theta$ is a random variable.
$\diamond$ Use a parametric optimization algorithm, i.e. a measurable function:

$$
\mathcal{A}: \mathcal{H} \times \mathbb{R}^{n} \times \Theta \rightarrow \mathbb{R}^{n}, \quad\left(\alpha, x^{(0)}, \theta\right) \mapsto \mathcal{A}\left(\alpha, x^{(0)}, \theta\right)
$$

## Example:

Regularized Inverse Problem:

$$
\ell(x, \lambda)=\frac{1}{2}\|A x-b\|^{2}+\lambda R(x), \quad \text { i.e. } \theta:=\lambda, \Theta=[0,1] .
$$

Concatenation of a fixed number of Preconditioned Gradient Descent steps:

$$
x^{(k+1)}=x^{(k)}-P \nabla \ell\left(x^{(k)}, \theta\right), \quad \text { i.e. } \alpha:=P, \mathcal{H}:=\mathbb{R}^{n \times n} .
$$

$\rightsquigarrow$ Learning boils down to hyperparameter optimization, i.e. how to choose $\alpha \in \mathcal{H}$.

## Quest for Theoretical Convergence Guarantees

Deterministic/Analytic Approach: Worst case performance

$$
\min _{\alpha \in \mathcal{H}} \sup _{\theta \in \Theta} \ell(\mathcal{A}(\alpha, \theta), \theta) .
$$

Only possible for certain classes of problems.

## Quest for Theoretical Convergence Guarantees

Deterministic/Analytic Approach: Worst case performance

$$
\min _{\alpha \in \mathcal{H}} \sup _{\theta \in \Theta} \ell(\mathcal{A}(\alpha, \theta), \theta) .
$$

Only possible for certain classes of problems.

Learning Based Approach: Expected case performance:
$\diamond$ Minimize the risk $\mathcal{R}(\alpha)$, defined as the expected loss:

$$
\min _{\alpha \in \mathcal{H}} \mathcal{R}(\alpha), \quad \mathcal{R}(\alpha):=\mathbb{E}[\ell(\mathcal{A}(\alpha, \mathfrak{S}), \mathfrak{S})]
$$

## Quest for Theoretical Convergence Guarantees

Deterministic/Analytic Approach: Worst case performance

$$
\min _{\alpha \in \mathcal{H}} \sup _{\theta \in \Theta} \ell(\mathcal{A}(\alpha, \theta), \theta) .
$$

Only possible for certain classes of problems.

Learning Based Approach: Expected case performance:
$\diamond$ Minimize the risk $\mathcal{R}(\alpha)$, defined as the expected loss:

$$
\min _{\alpha \in \mathcal{H}} \mathcal{R}(\alpha), \quad \mathcal{R}(\alpha):=\mathbb{E}[\ell(\mathcal{A}(\alpha, \mathfrak{S}), \mathfrak{S})] .
$$

This is intractable, since the distribution $\mathbb{P}_{\mathfrak{S}}$ is unknown.

## Quest for Theoretical Convergence Guarantees

Deterministic/Analytic Approach: Worst case performance

$$
\min _{\alpha \in \mathcal{H}} \sup _{\theta \in \Theta} \ell(\mathcal{A}(\alpha, \theta), \theta) \text {. }
$$

Only possible for certain classes of problems.

Learning Based Approach: Expected case performance:
$\diamond$ Minimize the risk $\mathcal{R}(\alpha)$, defined as the expected loss:

$$
\min _{\alpha \in \mathcal{H}} \mathcal{R}(\alpha), \quad \mathcal{R}(\alpha):=\mathbb{E}[\ell(\mathcal{A}(\alpha, \mathfrak{S}), \mathfrak{S})] .
$$

This is intractable, since the distribution $\mathbb{P}_{\mathfrak{E}}$ is unknown.
$\diamond$ Hence, resort to minimizing the empirical risk $\hat{\mathcal{R}}\left(\alpha, \mathfrak{D}_{N}\right)$ over some dataset $\mathfrak{D}_{N}:=\left\{\mathfrak{S}_{i}\right\}_{i=1}^{N}:$

$$
\min _{\alpha \in \mathcal{H}} \hat{\mathcal{R}}\left(\alpha, \mathfrak{D}_{N}\right), \quad \hat{\mathcal{R}}\left(\alpha, \mathfrak{D}_{N}\right):=\frac{1}{N} \sum_{i=1}^{N} \ell\left(\mathcal{A}\left(\alpha, \mathfrak{S}_{i}\right), \mathfrak{S}_{i}\right) .
$$

## Why do we need Generalization Guarantees?

Is the performance on $\hat{\mathcal{R}}$ representative for the overall performance $\mathcal{R}$ ?

## Why do we need Generalization Guarantees?

Is the performance on $\hat{\mathcal{R}}$ representative for the overall performance $\mathcal{R}$ ?
$\diamond$ Yes, if we have uniform generalization bounds, i.e. bounds of the form: $\forall \varepsilon>0$ :

$$
\mathbb{P}\left\{\mathcal{R}\left(\alpha^{*}\left(\mathfrak{D}_{N}\right)\right) \leq \inf _{\alpha \in \mathcal{H}} \hat{\mathcal{R}}\left(\alpha, \mathfrak{D}_{N}\right)+K(N, \alpha, \epsilon)\right\} \geq 1-\epsilon
$$

## Why do we need Generalization Guarantees?

Is the performance on $\hat{\mathcal{R}}$ representative for the overall performance $\mathcal{R}$ ?
$\diamond$ Yes, if we have uniform generalization bounds, i.e. bounds of the form: $\forall \varepsilon>0$ :

$$
\mathbb{P}\left\{\mathcal{R}\left(\alpha^{*}\left(\mathfrak{D}_{N}\right)\right) \leq \inf _{\alpha \in \mathcal{H}} \hat{\mathcal{R}}\left(\alpha, \mathfrak{D}_{N}\right)+K(N, \alpha, \epsilon)\right\} \geq 1-\epsilon
$$

## Why do we need Generalization Guarantees?

Is the performance on $\hat{\mathcal{R}}$ representative for the overall performance $\mathcal{R}$ ?
$\diamond$ Yes, if we have uniform generalization bounds, i.e. bounds of the form: $\forall \varepsilon>0$ :

$$
\mathbb{P}\left\{\mathcal{R}\left(\alpha^{*}\left(\mathfrak{D}_{N}\right)\right) \leq \inf _{\alpha \in \mathcal{H}} \hat{\mathcal{R}}\left(\alpha, \mathfrak{D}_{N}\right)+K(N, \alpha, \epsilon)\right\} \geq 1-\epsilon
$$

## Why do we need Generalization Guarantees?

Is the performance on $\hat{\mathcal{R}}$ representative for the overall performance $\mathcal{R}$ ?
$\diamond$ Yes, if we have uniform generalization bounds, i.e. bounds of the form: $\forall \varepsilon>0$ :

$$
\mathbb{P}\left\{\mathcal{R}\left(\alpha^{*}\left(\mathfrak{D}_{N}\right)\right) \leq \inf _{\alpha \in \mathcal{H}} \hat{\mathcal{R}}\left(\alpha, \mathfrak{D}_{N}\right)+K(N, \alpha, \epsilon)\right\} \geq 1-\epsilon .
$$

## Is the performance on $\hat{\mathcal{R}}$ representative for the overall performance $\mathcal{R}$ ?

$\diamond$ Yes, if we have uniform generalization bounds, i.e. bounds of the form: $\forall \varepsilon>0$ :

$$
\mathbb{P}\left\{\mathcal{R}\left(\alpha^{*}\left(\mathfrak{D}_{N}\right)\right) \leq \inf _{\alpha \in \mathcal{H}} \hat{\mathcal{R}}\left(\alpha, \mathfrak{D}_{N}\right)+K(N, \alpha, \epsilon)\right\} \geq 1-\epsilon .
$$

Such bounds are called PAC-bounds, which is an acronym for:


With high probability, the empirical risk is close to the true risk.

## PAC-Bayes extends this to the Bayes-risk:

Such bounds hold for posterior distributions $\mathbb{Q} \in \mathcal{M}\left(\mathbb{P}_{\mathfrak{F}}\right)$ :

$$
\mathbb{P}\left\{\mathbb{E}_{\mathbb{Q}^{*}\left(\mathscr{D}_{N}\right)}[\mathcal{R}] \leq \inf _{\mathbb{Q} \in \mathcal{M}\left(\mathbb{P}_{\mathfrak{S})}\right)} \mathbb{E}_{\mathbb{Q}}\left[\hat{\mathcal{R}}\left(\mathfrak{D}_{N}\right)\right]+K(\mathbb{Q}, N, \epsilon)\right\} \geq 1-\epsilon,
$$

where $\mathcal{M}\left(\mathbb{P}_{\mathfrak{H}}\right)$ denotes some class of (probability) measures on $\mathcal{H}$ that satisfy a certain property w.r.t. the prior distribution $\mathbb{P}_{\mathfrak{f}}$.

This is a naming convention! Not to be confused with prior and posterior in Bayesian analysis, which are linked by a likelihood.

## For good reviews of two long lines of work see...

PAC-Bayes [Alquier '21]


PAC-Bayesian Learning of Optimization Algorithms [Sucker, ©. 22]

[^0]$\diamond$ [Chen et al. '22]: "Learning to optimize: A primer and a benchmark", Journal of Machine Learning Research (2022), pp. 8562-8620.

## Towards a PAC-Bayes Theorem for Exponential Families

## A Form of the Donsker-Varadhan Variational Formulation

Lemma: Consider an exponential family $\left(\mathbb{Q}_{\lambda}\right)_{\lambda \in \Lambda}$ w.r.t. the prior $\mathbb{P}_{\mathfrak{H}}$, i.e. distributions of the form:

$$
\mathbb{Q}_{\lambda} \propto \exp (\langle\eta(\lambda), T\rangle) \cdot \mathbb{P}_{\mathfrak{j}}, \quad \lambda \in \Lambda
$$

and denote $c(\lambda):=\mathbb{E}_{\mathbb{P}_{\mathfrak{S}}}[\exp (\langle\eta(\lambda), T\rangle)]$. Then it holds:

$$
\log (c(\lambda))=\sup _{\mathbb{Q} \ll \mathbb{P}_{\mathfrak{H}}} \mathbb{E}_{\mathbb{Q}}[\langle\eta(\lambda), T\rangle]-D_{K L}\left(\mathbb{Q} \| \mathbb{P}_{\mathfrak{H}}\right)
$$

Furthermore, for every $\lambda \in \Lambda$, the supremum is attained at $\mathbb{Q}_{\lambda}$.

## Towards a PAC-Bayes Theorem for Exponential Families

## A Form of the Donsker-Varadhan Variational Formulation

Lemma: Consider an exponential family $\left(\mathbb{Q}_{\lambda}\right)_{\lambda \in \Lambda}$ w.r.t. the prior $\mathbb{P}_{\mathfrak{H}}$, i.e. distributions of the form:

$$
\mathbb{Q}_{\lambda} \propto \exp (\langle\eta(\lambda), T\rangle) \cdot \mathbb{P}_{\mathfrak{j}}, \quad \lambda \in \Lambda
$$

and denote $c(\lambda):=\mathbb{E}_{\mathbb{P}_{5}}[\exp (\langle\eta(\lambda), T\rangle)]$. Then it holds:

$$
\log (c(\lambda))=\sup _{\mathbb{Q} \ll \mathbb{P}_{\mathfrak{H}}} \mathbb{E}_{\mathbb{Q}}[\langle\eta(\lambda), T\rangle]-D_{K L}\left(\mathbb{Q} \| \mathbb{P}_{\mathfrak{H}}\right)
$$

Furthermore, for every $\lambda \in \Lambda$, the supremum is attained at $\mathbb{Q}_{\lambda}$.

## Towards a PAC-Bayes Theorem for Exponential Families

## A Form of the Donsker-Varadhan Variational Formulation

Lemma: Consider an exponential family $\left(\mathbb{Q}_{\lambda}\right)_{\lambda \in \Lambda}$ w.r.t. the prior $\mathbb{P}_{\mathfrak{H}}$, i.e. distributions of the form:

$$
\mathbb{Q}_{\lambda} \propto \exp (\langle\eta(\lambda), T\rangle) \cdot \mathbb{P}_{\mathfrak{H}}, \quad \lambda \in \Lambda
$$

and denote $c(\lambda):=\mathbb{E}_{\mathbb{P}_{\mathfrak{S}}}[\exp (\langle\eta(\lambda), T\rangle)]$. Then it holds:

$$
\log (c(\lambda))=\sup _{\mathbb{Q} \ll \mathbb{P}_{\mathfrak{H}}} \mathbb{E}_{\mathbb{Q}}[\langle\eta(\lambda), T\rangle]-D_{K L}\left(\mathbb{Q} \| \mathbb{P}_{\mathfrak{H}}\right)
$$

Furthermore, for every $\lambda \in \Lambda$, the supremum is attained at $\mathbb{Q}_{\lambda}$.

## Towards a PAC-Bayes Theorem for Exponential Families

## A Form of the Donsker-Varadhan Variational Formulation

Lemma: Consider an exponential family $\left(\mathbb{Q}_{\lambda}\right)_{\lambda \in \Lambda}$ w.r.t. the prior $\mathbb{P}_{\mathfrak{H}}$, i.e. distributions of the form:

$$
\mathbb{Q}_{\lambda} \propto \exp (\langle\eta(\lambda), T\rangle) \cdot \mathbb{P}_{\mathfrak{H}}, \quad \lambda \in \Lambda
$$

and denote $c(\lambda):=\mathbb{E}_{\mathbb{P}_{\mathfrak{S}}}[\exp (\langle\eta(\lambda), T\rangle)]$. Then it holds:

$$
\log (c(\lambda))=\sup _{\mathbb{Q} \ll \mathbb{P}_{\mathfrak{H}}} \mathbb{E}_{\mathbb{Q}}[\langle\eta(\lambda), T\rangle]-D_{K L}\left(\mathbb{Q} \| \mathbb{P}_{\mathfrak{H}}\right)
$$

Furthermore, for every $\lambda \in \Lambda$, the supremum is attained at $\mathbb{Q}_{\lambda}$.

## Towards a PAC-Bayes Theorem for Exponential Families

Theorem: If $\mathbb{E}_{\mathfrak{D}_{N}}[c(\lambda)] \leq 1$, then for all $\varepsilon>0$ :

$$
\mathbb{P}_{\mathfrak{D}_{N}}\left\{\forall \lambda \in \Lambda: \forall \mathbb{Q} \ll \mathbb{P}_{\mathfrak{5}}: \mathbb{E}_{\mathbb{Q}}[\langle\eta(\lambda), T\rangle] \leq D_{\mathrm{KL}}\left(\mathbb{Q} \| \mathbb{P}_{\mathfrak{5}}\right)+\log (|\Lambda| / \varepsilon)\right\} \geq 1-\varepsilon
$$

Theorem: If $\mathbb{E}_{\mathfrak{D}_{N}}[c(\lambda)] \leq 1$, then for all $\varepsilon>0$ :

$$
\mathbb{P}_{\mathfrak{D}_{N}}\left\{\forall \lambda \in \Lambda: \forall \mathbb{Q} \ll \mathbb{P}_{\mathfrak{5}}: \mathbb{E}_{\mathbb{Q}}[\langle\eta(\lambda), T\rangle] \leq D_{\mathrm{KL}}\left(\mathbb{Q} \| \mathbb{P}_{\mathfrak{5}}\right)+\log (|\Lambda| / \varepsilon)\right\} \geq 1-\varepsilon
$$

## Sketch of Proof.

Use Markov's inequality

$$
\mathbb{P}_{\mathfrak{P}_{N}}\{c(\lambda) \geq \exp (s)\} \leq \frac{\mathbb{E}_{\mathfrak{Q}_{N}}[c(\lambda)]}{\exp (s)} \leq 1 / \exp (s)=: 1 / s^{\prime} .
$$

Theorem: If $\mathbb{E}_{\mathfrak{D}_{N}}[c(\lambda)] \leq 1$, then for all $\varepsilon>0$ :

$$
\mathbb{P}_{\mathfrak{D}_{N}}\left\{\forall \lambda \in \Lambda: \forall \mathbb{Q} \ll \mathbb{P}_{\mathfrak{S}}: \mathbb{E}_{\mathbb{Q}}[\langle\eta(\lambda), T\rangle] \leq D_{\mathrm{KL}}\left(\mathbb{Q} \| \mathbb{P}_{\mathfrak{5}}\right)+\log (|\Lambda| / \varepsilon)\right\} \geq 1-\varepsilon
$$

## Sketch of Proof.

$\diamond$ Use Markov's inequality

$$
\mathbb{P}_{\mathfrak{P}_{N}}\{c(\lambda) \geq \exp (s)\} \leq \frac{\mathbb{E}_{\mathfrak{D}_{N}}[c(\lambda)]}{\exp (s)} \leq 1 / \exp (s)=: 1 / s^{\prime} .
$$

$\diamond$ Union-bound argument: (use covering argument for compact continuous $\Lambda$ )

$$
\mathbb{P}_{\mathfrak{O}_{N}}\left\{\sup _{\lambda \in \Lambda} c(\lambda)>s^{\prime}\right\}=\mathbb{P}_{\mathfrak{D}_{N}}\left\{\bigcup_{\lambda \in \Lambda}\left\{c(\lambda)>s^{\prime}\right\}\right\} \leq \sum_{\lambda \in \Lambda} \mathbb{P}_{\mathfrak{D}_{N}}\left\{\left\{c(\lambda)>s^{\prime}\right\}\right\} \leq|\Lambda| / s^{\prime}=: \varepsilon
$$

## Towards a PAC-Bayes Theorem for Exponential Families

Theorem: If $\mathbb{E}_{\mathfrak{D}_{N}}[c(\lambda)] \leq 1$, then for all $\varepsilon>0$ :

$$
\mathbb{P}_{\mathfrak{D}_{N}}\left\{\forall \lambda \in \Lambda: \forall \mathbb{Q} \ll \mathbb{P}_{\mathfrak{S}}: \mathbb{E}_{\mathbb{Q}}[\langle\eta(\lambda), T\rangle] \leq D_{\mathrm{KL}}\left(\mathbb{Q} \| \mathbb{P}_{\mathfrak{5}}\right)+\log (|\Lambda| / \varepsilon)\right\} \geq 1-\varepsilon
$$

## Sketch of Proof.

$\diamond$ Use Markov's inequality

$$
\mathbb{P}_{\mathfrak{D}_{N}}\{c(\lambda) \geq \exp (s)\} \leq \frac{\mathbb{E}_{\mathfrak{D}_{N}}[c(\lambda)]}{\exp (s)} \leq 1 / \exp (s)=: 1 / s^{\prime} .
$$

$\diamond$ Union-bound argument: (use covering argument for compact continuous $\Lambda$ )

$$
\mathbb{P}_{\mathcal{O}_{N}}\left\{\sup _{\lambda \in \Lambda} c(\lambda)>s^{\prime}\right\}=\mathbb{P}_{\mathcal{O}_{N}}\left\{\bigcup_{\lambda \in \Lambda}\left\{c(\lambda)>s^{\prime}\right\}\right\} \leq \sum_{\lambda \in \Lambda} \mathbb{P}_{\mathfrak{O}_{N}}\left\{\left\{c(\lambda)>s^{\prime}\right\}\right\} \leq|\Lambda| / s^{\prime}=: \varepsilon
$$

$\diamond$ Apply Donsker-Varadhan variational formulation in

$$
\mathbb{P}_{\mathfrak{D}_{N}}\left\{\sup _{\lambda \in \Lambda} \log (c(\lambda)) \leq \log (|\Lambda| / \varepsilon)\right\} \geq 1-\varepsilon .
$$

Specify $\eta$ and $T$ accordingly to construct a PAC-Bayesian generalization bound:

$$
\eta(\lambda)=\left(\eta_{1}(\lambda), \eta^{\prime}(\lambda)\right)
$$

and

$$
T\left(\alpha, \mathfrak{D}_{N}\right)=\left(\mathcal{R}(\alpha)-\hat{\mathcal{R}}\left(\alpha, \mathfrak{D}_{N}\right), T^{\prime}\left(\alpha, \mathfrak{D}_{N}\right)\right)
$$

Specify $\eta$ and $T$ accordingly to construct a PAC-Bayesian generalization bound:

$$
\eta(\lambda)=\left(\eta_{1}(\lambda), \eta^{\prime}(\lambda)\right)
$$

and

$$
T\left(\alpha, \mathfrak{D}_{N}\right)=\left(\mathcal{R}(\alpha)-\hat{\mathcal{R}}\left(\alpha, \mathfrak{D}_{N}\right), T^{\prime}\left(\alpha, \mathfrak{D}_{N}\right)\right) .
$$

This provides a bound of the following form:
Theorem: Under mild assumptions, it holds for $\epsilon>0$ :

$$
\mathbb{P}_{\mathfrak{D}_{N}}\left\{\forall \lambda \in \Lambda, \forall \mathbb{Q} \ll \mathbb{P}_{\mathfrak{H}}: \mathbb{E}_{\mathbb{Q}}[\mathcal{R}] \leq \mathbb{E}_{\mathbb{Q}}[\hat{\mathcal{R}}]+G(N, \lambda, \mathbb{Q}, \epsilon)\right\} \geq 1-\epsilon .
$$

## Note:

(i) By the definition of the risk and the algorithm, this bound gives a guarantee for the function value of the algorithm's output.
(ii) This is a statement about relative values, not absolute ones.

Specify $\eta$ and $T$ accordingly to construct a PAC-Bayesian generalization bound:

$$
\eta(\lambda)=\left(\eta_{1}(\lambda), \eta^{\prime}(\lambda)\right)
$$

and

$$
T\left(\alpha, \mathfrak{D}_{N}\right)=\left(\mathcal{R}(\alpha)-\hat{\mathcal{R}}\left(\alpha, \mathfrak{D}_{N}\right), T^{\prime}\left(\alpha, \mathfrak{D}_{N}\right)\right) .
$$

This provides a bound of the following form:
Theorem: Under mild assumptions, it holds for $\epsilon>0$ :

$$
\mathbb{P}_{\mathfrak{D}_{N}}\left\{\forall \lambda \in \Lambda, \forall \mathbb{Q} \ll \mathbb{P}_{\mathfrak{H}}: \mathbb{E}_{\mathbb{Q}}[\mathcal{R}] \leq \mathbb{E}_{\mathbb{Q}}[\hat{\mathcal{R}}]+G(N, \lambda, \mathbb{Q}, \epsilon)\right\} \geq 1-\epsilon .
$$

## Note:

(i) By the definition of the risk and the algorithm, this bound gives a guarantee for the function value of the algorithm's output.
(ii) This is a statement about relative values, not absolute ones.

Since supremum is attained at $\mathbb{Q}_{\lambda}$, learning can be phrased as an optimization in $\lambda$ (possibly very low-dimensional).

## Simple Case Study: Gradient Descent

## Issue: A bad performance on a single problem dominates the average.

$\diamond$ Sometimes, analytic worst-case bounds are sharp.
$\diamond$ Gradient Descent on quadratics diverges for $\alpha>2 / L$.
$\diamond$ Trying to learn the step size (without this bound) yields an extremely large loss for $\alpha>2 / L$, which dominates the cost of the "average performance" (the empirical risk). Therefore, learnable step sizes obey the deterministic step size rule $\alpha \in(0,2 / L)$.


In this case, the analytically best known step size is recoverd by learning.

## Trade-Off Guarantees and Speed

However this might correspond to unlikely special/extreme cases:

## We develop a variant that allows to Trade-Off Guarantees and Speed.

## Trade-Off Guarantees and Speed

However this might correspond to unlikely special/extreme cases:

We develop a variant that allows to Trade-Off Guarantees and Speed.

The algorithm may diverge (for extreme cases), if this happens in rare cases and the Trade-Off can be controlled.

## Trade-Off Guarantees and Speed

However this might correspond to unlikely special/extreme cases:

We develop a variant that allows to Trade-Off Guarantees and Speed.

The algorithm may diverge (for extreme cases), if this happens in rare cases and the Trade-Off can be controlled.

Account for likelihood of e.g. the worst-case.

## Trade-Off Guarantees and Speed

However this might correspond to unlikely special/extreme cases:

We develop a variant that allows to Trade-Off Guarantees and Speed.

The algorithm may diverge (for extreme cases), if this happens in rare cases and the Trade-Off can be controlled.

Account for likelihood of e.g. the worst-case.
Encode properties of the algorithm in the convergence set $\mathrm{C} \subset \mathcal{H} \times \Theta$, e.g.,

$$
\mathrm{C}_{\alpha}:=\left\{\theta \in \Theta: \ell(\mathcal{A}(\alpha, \theta), \theta) \leq \ell\left(x^{(0)}, \theta\right)\right\}
$$

## Trade-Off Guarantees and Speed

However this might correspond to unlikely special/extreme cases:

We develop a variant that allows to Trade-Off Guarantees and Speed.

The algorithm may diverge (for extreme cases), if this happens in rare cases and the Trade-Off can be controlled.

Account for likelihood of e.g. the worst-case.
Encode properties of the algorithm in the convergence set $\mathrm{C} \subset \mathcal{H} \times \Theta$, e.g.,

$$
\mathrm{C}_{\alpha}:=\left\{\theta \in \Theta: \ell(\mathcal{A}(\alpha, \theta), \theta) \leq \ell\left(x^{(0)}, \theta\right)\right\}
$$

Condition on it to get the convergence risk $\mathcal{R}_{c}: \mathcal{H} \rightarrow \mathbb{R}_{\geq 0}$ :

$$
\mathcal{R}_{c}(\alpha):=\mathbb{E}\left[\ell(\mathcal{A}(\alpha, \mathfrak{S}), \mathfrak{S}) \mid \mathrm{C}_{\alpha}\right]
$$

## Trade-Off Guarantees and Speed

However this might correspond to unlikely special/extreme cases:

We develop a variant that allows to Trade-Off Guarantees and Speed.

The algorithm may diverge (for extreme cases), if this happens in rare cases and the Trade-Off can be controlled.

Account for likelihood of e.g. the worst-case.
Encode properties of the algorithm in the convergence set $\mathrm{C} \subset \mathcal{H} \times \Theta$, e.g.,

$$
\mathrm{C}_{\alpha}:=\left\{\theta \in \Theta: \ell(\mathcal{A}(\alpha, \theta), \theta) \leq \ell\left(x^{(0)}, \theta\right)\right\}
$$

Condition on it to get the convergence risk $\mathcal{R}_{c}: \mathcal{H} \rightarrow \mathbb{R}_{\geq 0}$ :

$$
\mathcal{R}_{c}(\alpha):=\mathbb{E}\left[\ell(\mathcal{A}(\alpha, \mathfrak{S}), \mathfrak{S}) \mid \mathrm{C}_{\alpha}\right]
$$

$\diamond$ Guarantees in form of the convergence probability $\mathbb{P}_{\mathfrak{S}}\left[\mathrm{C}_{\alpha}\right]$ instead of convergence for every sample.

## Trade-Off Guarantees and Speed

Applying the same machinery again yields the following generalization:

Theorem: Under mild assumptions, it holds for $\varepsilon>0$ :

$$
\mathbb{P}_{\mathfrak{D}_{N}}\left\{\forall \lambda \in \Lambda, \forall \mathbb{Q} \ll \mathbb{P}_{\mathfrak{H}}: \mathbb{E}_{\mathbb{Q}}\left[\mathcal{R}_{c}\right] \leq \mathbb{E}_{\mathbb{Q}}\left[\hat{\mathcal{R}}_{c}\right]+G(N, \lambda, \mathbb{Q}, \varepsilon)\right\} \geq 1-\varepsilon .
$$

## Trade-Off Guarantees and Speed

Applying the same machinery again yields the following generalization:

Theorem: Under mild assumptions, it holds for $\varepsilon>0$ :

$$
\mathbb{P}_{\mathfrak{D}_{N}}\left\{\forall \lambda \in \Lambda, \forall \mathbb{Q} \ll \mathbb{P}_{\mathfrak{H}}: \mathbb{E}_{\mathbb{Q}}\left[\mathcal{R}_{c}\right] \leq \mathbb{E}_{\mathbb{Q}}\left[\hat{\mathcal{R}}_{c}\right]+G(N, \lambda, \mathbb{Q}, \varepsilon)\right\} \geq 1-\varepsilon
$$

$\diamond$ Learning must achieve that $\mathcal{A}$ converges for "sufficiently many problems" (according to the convergence probability).

Therefore, the algorithm can focus on quickly solving the remaining problems.

## Trade-Off Guarantees and Speed

Applying the same machinery again yields the following generalization:

Theorem: Under mild assumptions, it holds for $\varepsilon>0$ :

$$
\mathbb{P}_{\mathfrak{D}_{N}}\left\{\forall \lambda \in \Lambda, \forall \mathbb{Q} \ll \mathbb{P}_{\mathfrak{H}}: \mathbb{E}_{\mathbb{Q}}\left[\mathcal{R}_{c}\right] \leq \mathbb{E}_{\mathbb{Q}}\left[\hat{\mathcal{R}}_{c}\right]+G(N, \lambda, \mathbb{Q}, \varepsilon)\right\} \geq 1-\varepsilon
$$

$\diamond$ Learning must achieve that $\mathcal{A}$ converges for "sufficiently many problems" (according to the convergence probability).

Therefore, the algorithm can focus on quickly solving the remaining problems.
Example Statement: With high probability, the algorithm that is trained to optimize $95 \%$ of all problems in $\mathfrak{D}_{N}$ quickly, will optimize $95 \%$ of all problems quickly.

The Whole Training Process


1) Find a "trainable" initialization by following another algorithm.

2) Run a specifically constrained sampling procedure.

3) Find a point inside the constraint with small empirical risk.

4) Find $\lambda^{*}$ and perform a reweighting based on closed-form of the posterior.


Training a 2-layer neural network with ReLUactivations ...


... to perform regression.


Cumulative Time to Solve the Test Set


## Learning gets faster...



Iteration 9



Iteration 12


Iteration 6


Iteration 15


## PAC-Bayes

Learning-to-Optimize


## PAC-Bayesian Learning of Optimization Algorithms [Sucker, 0. 22]

Breaking the barrier of worst-case estimates

$$
\min _{\alpha \in \mathcal{H}} \mathcal{R}(\alpha), \quad \mathcal{R}(\alpha):=\mathbb{E}[\ell(\mathcal{A}(\alpha, \mathfrak{S}), \mathfrak{S})]
$$

$\diamond$ by learning spezialized optimization algorithms

$$
\min _{\alpha \in \mathcal{H}} \hat{\mathcal{R}}\left(\alpha, \mathfrak{D}_{N}\right), \quad \hat{\mathcal{R}}\left(\alpha, \mathfrak{D}_{N}\right):=\frac{1}{N} \sum_{i=1}^{N} \ell\left(\mathcal{A}\left(\alpha, \mathfrak{S}_{i}\right), \mathfrak{S}_{i}\right) .
$$

$\diamond$ with theoretical guarantees via PAC-Bayes generalization bounds:

$$
\mathbb{P}_{\mathfrak{D}_{N}}\left\{\forall \lambda \in \Lambda, \forall \mathbb{Q} \ll \mathbb{P}_{\mathfrak{H}}: \mathbb{E}_{\mathbb{Q}}[\mathcal{R}] \leq \mathbb{E}_{\mathbb{Q}}[\hat{\mathcal{R}}]+G(N, \lambda, \mathbb{Q}, \epsilon)\right\} \geq 1-\epsilon
$$


[^0]:    $\diamond$ [Alquier '21]: "User-friendly introduction to PAC-Bayes bounds", arXiv:2110.11216 (2021).

